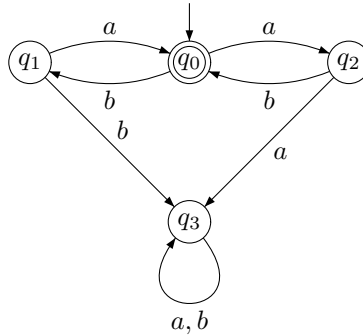


**ASSIGNMENT FOR THE COURSE
THEORETICAL COMPUTER SCIENCE
THEORY OF ALGORITHMS**

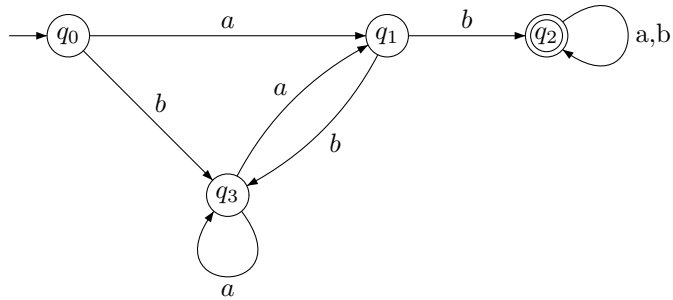
- (1) Prove that $\{ab, aba\}^* = \{\varepsilon\} \cup \{a\}\{ba, baa\}^*\{b, ba\}$ (i.e. prove 2 inclusions: “ \subseteq ” and “ \supseteq ”).
- (2) Build a DFA which recognizes the language $\{0, 1\}^*$ in alphabet $\mathcal{A} = \{0, 1, 2\}$.
- (3) Build a DFA which recognizes

$L = \{w \mid w \in \{0, 1\}^* \text{ and the difference between the number of 0's and 1's in } w \text{ is even}$
in $\mathcal{A} = \{0, 1\}$.

- (4) Build a DFA for $L = \{w \mid \text{in } w \text{ before } a \text{ always comes } b\}$ in $\mathcal{A} = \{a, b\}$.
- (5) Describe the language recognized by the following DFA:



- (6) Describe the language recognized by the following NFA:



- (7) Using the algorithm considered in the lecture which allows one to build a DFA from a NFA, build a DFA for the language from Exercise (6).
- (8) Construct a regular grammar generating $L = \{w \mid w \neq \varepsilon \text{ and the number of 1's divides by 3}\}$.
- (9) Prove that every finite language not containing ε is generated by a regular grammar.
- (10) Prove that $L = \{w \in \{a, b\}^+ \mid w \text{ has equal number of occurrences of } a \text{ and } b\}$ is not regular.
- (11) Prove that $L = \{ww^R \mid w \in \{a, b\}^+\}$ is not regular.
- (12) Build a PDA recognizing $L = \{a^n b^n \mid n \in \omega\}$.

- (13) Describe the language generated by the following CFG $\Gamma = \langle V, T, P, S \rangle$, where $V = \{S, A, B\}$, $T = \{a, b\}$, $P = \{S \rightarrow aB, S \rightarrow bA, A \rightarrow a, A \rightarrow aS, A \rightarrow BAA, B \rightarrow b, B \rightarrow bS, B \rightarrow ABB\}$.

- (14) Construct a CFG which generates

$$L = \{xc^n \mid n \in \omega, x \in \{a, b\}^* \text{ and the number of } a\text{'s is } n \text{ or the number of } b\text{'s is } n\}$$

- (15) Provide an example of context-free languages L_1, L_2 such that $L_1 \setminus L_2$ is not context-free.

- (16) Prove that $L = \{ww \mid w \in \{0, 1\}^*\}$ is not context-free.

- (17) Prove that $L = \{a^i b^j c^k \mid i < j < k\}$ is not context-free.

- (18) Prove that $f(x) = x!$ (where $0! = 1$) is primitive recursive.

- (19) Prove that the function $rest(x, y)$ which outputs the remainder of division x by y (where $rest(x, 0) = x$) is primitive recursive.

- (20) Prove that $f(x) = \lfloor \sqrt{x} \rfloor$ is primitive recursive.

- (21) Describe the function which is the result of application of the minimization operator to $g(x, y, z) = |zy - x|$.

- (22) Prove that the function f such that $f(n)$ is the n -th Fibonacci number, is primitive recursive.

- (23) Construct a Turing machine which performs transposition: $01^x q_1 01^y 0 \Rightarrow 01^y q_0 01^x 0$ *without building new cells to the left and to the right end of the tape*.

- (24) Construct a Turing machine which performs duplication: $q_1 01^x 0 \Rightarrow q_0 01^x 01^x 0$.

- (25) Construct a Turing machine which correctly computes the function $f(x, y) = x - y$ (in particular, if $x < y$ the machine does not halt).

A set A is computable if its characteristic function χ_A is computable (where $\chi_A(x) = 1$ if $x \in A$, $\chi_A = 0$ if $x \notin A$).

- (26) Prove that if $f(x)$ is computable then its graph Γ_f is a computable set.

- (27) Prove that if $f(x)$ is a monotone computable function then its range is a computable set.

- (28) Prove that if $f(x, y)$ is computable and $g(x) = \mu y[f(x, y) = 0]$ then the graph of g Γ_g is computable.

- (29) Prove that the set $A = \{x \in \omega \mid x \text{ is a perfect number}\}$ is computable (a number x is perfect if the sum of all its divisors that are less than x equals x).

- (30) Prove that it is impossible to get the functions $s(x) = x + 1$ and $f(x) = 2x$ from the basic functions o and I_m^n using the operators of composition and primitive recursion.

- (31) Let $e = a_0, a_1 a_2 a_3 \dots$ be a representation of the number e (Euler's number 2, 718281...) as an infinite decimal fraction. Prove that $f(n) = a_n$ is computable.

- (32) Prove that the problem 2-SAT is in \mathcal{P} .

- (33) Let $L_1 \in \mathcal{P}$, L_2 be \mathcal{NP} -complete and $L_3 \notin \mathcal{NP}$. Characterize the languages $L_1 \cap L_2$, $L_1 \cup L_2$, $L_2 c L_3$ (where c is a symbol not in L_2 or L_3), $\overline{L_3}$ using the following expressions:

- (a) definitely in \mathcal{P} ;
- (b) definitely in \mathcal{NP} (but perhaps not in \mathcal{P} and perhaps not \mathcal{NP} -complete);
- (c) definitely \mathcal{NP} -complete;
- (d) definitely not in \mathcal{NP} .

Explain your choice.

- (34) Decide whether \mathcal{P} is closed under union, intersection, concatenation and star. Give a proof of your choice.
- (35) Decide whether \mathcal{NP} is closed under union, intersection, concatenation and star. Give a proof of your choice.
- (36) Build a nondeterministic Turing machine which recognizes $L = \{ww \mid w \in \{a, b\}^*\}$.