

6.0/4.0 VU Formale Methoden der Informatik				
185.291		SS 2011	28 October 2011	
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1.) Consider the following problem:

IMPERFECT-COMPRESSION(IC)

INSTANCE: A program (i.e. a source code) Π such that Π takes one string as input and outputs a string. It is guaranteed that Π terminates on any input string.

QUESTION: Does there exist an input string I for Π such that $|\Pi(I)| > |I|$. Here $|J|$ denotes the length of a string J , and $\Pi(J)$ is the string returned by Π on input string J .

Prove that the problem **IC** is semi-decidable. For this, describe a procedure that shows the semi-decidability of the problem (i.e. a semi-decision procedure for **IC**) and argue that it is correct. **(15 points)**

2.) Let φ be the formula $(x \rightarrow (y \rightarrow z)) \rightarrow ((x \wedge y) \rightarrow z)$.

(a) Use the Tseitin translation and compute $\hat{\delta}(\varphi)$. (Hint: It is allowed to avoid the labels for atoms; use the atoms instead. Moreover, use ℓ_φ as the label for φ .) **(5 points)**

(b) Try to derive the empty clause \square from

$$\left(\bigwedge_{D \in \hat{\delta}(\varphi)} D \right) \wedge \neg \ell_\varphi$$

by resolution.

(4 points)

(c) Answer the following questions and explain in detail.

- i. Is $\bigwedge_{D \in \hat{\delta}(\varphi)} D$ satisfiable? If so then provide a model.
- ii. Is $\bigwedge_{D \in \hat{\delta}(\varphi)} D$ valid? If not then provide a counterexample.
- iii. Is $\left(\bigwedge_{D \in \hat{\delta}(\varphi)} D \right) \rightarrow \ell_\varphi$ valid?
- iv. Is φ valid?

(6 points)

3.) Show that the following correctness assertion is totally correct. Describe the function computed by the program if we consider a as its input and c as its output.

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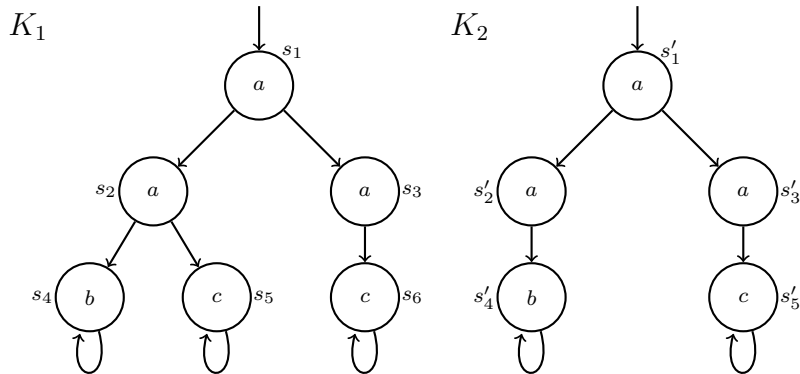
{ 1: a ≥ 0 }
b ← 1;
c ← 0;
{ Inv: b = (c + 1)3 ∧ 0 ≤ c3 ≤ a }
while b ≤ a do
  d ← 3 * c + 6;
  c ← c + 1;
  b ← b + c * d + 1
od
{ 2: c3 ≤ a < (c + 1)3 }

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(15 points)

4.) Simulation and Bisimulation

- (a) Let K_1 and K_2 be the two Kripke structures given below. Check which of the relations $K_1 \leq K_2$, $K_1 \geq K_2$, $K_1 \equiv K_2$ hold on K_1 and K_2 . Justify your answer.



(5 points)

- (b) Show that simulation is a transitive relation: Given any 3 Kripke structures $K_1 = \{S_1, R_1, L_1\}$, $K_2 = \{S_2, R_2, L_2\}$ and $K_3 = \{S_3, R_3, L_3\}$ such that $K_1 \leq K_2$ and $K_2 \leq K_3$, it holds that $K_1 \leq K_3$. (5 points)
- (c) Given the Kripke structures $K_1 = (S_1, R_1, L_1)$, $K_2 = (S_2, R_2, L_2)$, $K_3 = (S_3, R_3, L_3)$, the simulation relation $H_1 \subseteq S_1 \times S_2$ from K_1 to K_2 and the simulation relation $H_2 \subseteq S_2 \times S_3$ from K_2 to K_3 , state an algorithm which computes a simulation H_3 from K_1 to K_3 . (5 points)