VU Logic and Computation Exercises – Computability Theory

December 19, 2014

Preliminary information

- I expect *individual* solutions (no group work)
- The exercises should be written in Latex (preferred) or Word
- The (fix!) deadlines for the submission are:
 - January 9th 2015 12:00 *am*, if you want individual feedback on January 15th.
 - January 14th 2015 23:59, if you do not want individual feedback on January 15th.

No deadline extensions possible.

• The exercises should be sent via email (.pdf format) to

agata@logic.at

and the file should be named

*Computability.pdf

(where * is to be instantiated with your surname)

- An acknowledgement of receipt will be sent after the submission (no acknowledgement means no receipt)
- Suggestion: do not try to solve the exercises short before the deadline (and do not ask for any explanation on the exercises after January 7th, 2015 (12:00 am))

Exercises

Notation: We use Φ_k to denote the k-th computable function.

- 1. Lambda calculus
 - Exhibit a lambda term which encodes the function $f:\mathcal{N}\to\mathcal{N}$ s.t.

$$f(x) = 1 + \dots (x - 1) + x$$

- In arithmetic non every function has a fixpoint. Why, instead, every lambda term has a fixpoint?
- 2. Prove or disprove the following statements:
 - every countable union of recursive sets is recursive
 - every countable union of r.e. sets is r.e.
- 3. Is there is an index $p \in \mathcal{N}$ such that $\Phi_p \simeq \Phi_{p+1}$? Motivate your answer
- 4. Is the following function

$$g(x,y) = \begin{cases} 1 & \text{if } \forall z \ (\Phi_x(z) \downarrow \land \Phi_y(z) \downarrow \Rightarrow \ \Phi_x(z) = \Phi_y(z)) \\ \uparrow & \text{otherwise} \end{cases}$$

computable? Motivate your answer